

CS 3313:

Foundations of Computing

Lab 1: Proof Techniques

<http://gw-cs3313.github.io>

Outline



- Math preliminaries – this should all be review
- Proof techniques
- Exercises

Sets and set operations

- Set: Collection of non-repeating items $S = \{1, a, \{c, d\}\}$, $|S| = 3$

Sets and set operations

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 - \mathbb{Z} - set of integers
 - \mathbb{Z}^+ - set of positive integers
 - \mathbb{N} - natural numbers
 - \mathbb{R} - Reals
 - \emptyset or $\{\}$ – empty set

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 - \emptyset or $\{\}$ – empty set
- Set relations:
 - Membership: $5 \in \mathbb{Z}$, $3.1 \notin \mathbb{Z}$, $\{c, d\} \in \{1, a, \{c, d\}\}$
 - Subset: $\{1, 2\} \subset \{1, 2, 3\}$
 - Union: $A \cup B$ Intersection: $A \cap B$ Complement: \bar{A}
 - De Morgan's Laws: $\overline{A \cup B} = \bar{A} \cap \bar{B}$, $\overline{A \cap B} = \bar{A} \cup \bar{B}$
 - Cartesian product: If $A = \{1, 2, 3\}$, $B = \{a, b\}$
 $A \times B = \{(1, a), (1, b), (2, a), (2, b), (3, a), (3, b)\}$

Strings

- Alphabet Σ : set of symbols
 - Ex: $\Sigma = \{a, b\}$, $\Sigma = \{0,1\}$
- String: finite sequence of symbols from Σ ,
 - ex: $v = aba$ and $w = abaaa$
 - ex: $v = 001$ and $w = 11001$
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- Operations on strings:
 - Concatenation: $vw = abaabaaa$
 - Reverse: $w^R = aaaba$
 - Repetition: $v^2 = abaaba$ and $v^0 = \lambda$
- Length of a string: $|v| = 3$ and $|\lambda| = 0$

Languages

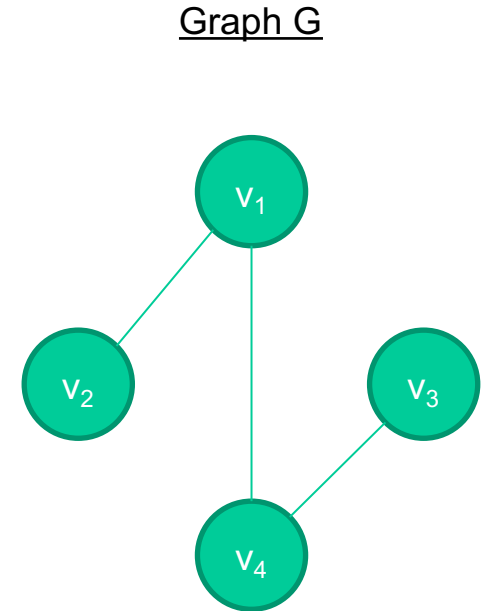
- Language L: Set of strings
- Σ^* = set of all strings formed by concatenating zero or more symbols in Σ
 - Ex: if $\Sigma = \{0,1\}$ then $\Sigma^* = \{\text{all binary strings, including empty string}\}$
- **A language is any subset of Σ^***

Examples: $L_1 = \{ a^n b^n : n \geq 0 \}$ and $L_2 = \{ ab, aa \}$


- A string in a language is also called a sentence of the language

Graphs

- A graph G consists of a
 - vertex set $V(G) = \{v_1, v_2, \dots\}$ and
 - edge set $E(G) \subset \{(x, y) | x, y \in V(G)\}$
i.e., an edge connects a pair of vertices
- Directed vs. undirected graphs
- Degree of a vertex – number of edges coming out of the vertex
e.g. $\deg(v_1) = 2$



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Proofs

This class will involve a lot of proofs.

General proof procedure:

1. Understand the statement – without the math lingo
2. Build up an intuition – in English or by picture, work through examples
 - This gets easier with practice
3. Construct proof
 - This part is procedural
 - Use facts and theorems you already know
 - Proof techniques will guide you

Writing proofs

- Be concise – no multi-paragraph explanations
- Be precise – use mathematic notation and logical reasoning
- Follow proof techniques – this will give you a structure for the proof

Proof techniques

- Direct proof
- Proof by contradiction
- Proof by induction

Direct Proof

- Produce a chain of logically sound deductions that justify the expected conclusion

Example

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- Proof: Important, there are two directions to prove!
 1. Suppose $x \in \overline{A \cup B}$, then $x \in \bar{A} \cap \bar{B}$
If $x \in \overline{A \cup B}$, then $x \notin A \cup B$ (by definition of complement)
So, $x \notin A$ and $x \notin B$ (by definition of union)
Thus, $x \in \bar{A}$ and $x \in \bar{B}$ (by definition of complement)
Therefore, $x \in \bar{A} \cap \bar{B}$ (by definition of intersection)

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Therefore, $x \in \bar{A} \cap \bar{B}$ (by definition of intersection)
 2. Suppose $x \in \bar{A} \cap \bar{B}$, then $x \in \overline{A \cup B}$
If $x \in \bar{A} \cap \bar{B}$, then $x \in \bar{A}$ and $x \in \bar{B}$ (by definition of intersection)
So, $x \notin A$ and $x \notin B$ (by definition of complement)
thus, $x \notin A \cup B$ (by definition of union)
implying that $x \in \overline{A \cup B}$ (by definition of complement)

Proof by contradiction

Proof Outline:

1. Assume the opposite of what you want to try to prove
2. Show that it leads to a contradiction
3. Thus, the original assumption must be false

Example

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- Intuition: The product of two odd numbers has no factors of 2, so it will be odd
- Proof:
 1. Assume there exists an even n s.t. n^2 is odd (very important to get the negation of the statement correct!)
Then, $n=2m$ for some integer m (by definition of even)
So, $n^2 = 4m^2 = 2(2m^2)$ which is even

Contradiction!!!

Proof by induction

Proof Outline:

1. **Base case:** Verify that statement holds for base case (e.g., true for $i=1$)
2. **Inductive hypothesis:** Assume that if the statement holds for $i=n$ for some value n
3. **Induction step:** Prove that the statement holds for $i=n+1$

Why this works:

$P(1)$ is true implies $P(2)$ is true

$P(2)$ is true implies $P(3)$ is true

...

$P(n-1)$ is true implies $P(n)$ is true

Therefore, $P(n)$ is true

Example

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- Proof:


Base case: $n = 1 - 1 = (1+1) * 1/2 = 2/2 = 1$

Hypothesis: Assume $\sum_{i=1}^k i = \frac{(k+1)k}{2}$ for some k

Induction: Show that $\sum_{i=1}^{k+1} i = \frac{((k+1)+1)(k+1)}{2}$

- $$\begin{aligned} \sum_{i=1}^{k+1} i &= k + 1 + \sum_{i=1}^k i = (k + 1) + \frac{(k+1)(k)}{2} \text{ (by hypothesis)} \\ &= \frac{k^2 + 3k + 2}{2} = \frac{(k+2)(k+1)}{2} \end{aligned}$$

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Exercises

- Prove each of the following statements
- Work in groups – make sure you put down all the names
- Scan your solutions and submit on GradeScope by end of day

Exercises

1. Prove that in any graph G , the sum of degrees of the nodes of G is an even number
2. Prove that $\sqrt{2}$ is irrational
3. Prove that

$$1^2 + 2^2 + \dots + n^2 = \sum_{i=1}^n i^2 = \frac{n(n+1)(2n+1)}{6}$$