

Foundations of Computing

Lecture 11

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February 20, 2024

- 1 Lecture 10 Review
- 2 The CFG Pumping Lemma
- 3 Midterm Review

Lecture 10 Review

- CFG \iff PDA
 - Construct PDA from CFG
 - Construct CFG from PDA
- CFG Pumping Lemma

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Today

- Midterm review

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The CFG Pumping Lemma

Theorem

If L is a CFL, then there exists a pumping length p s.t. for any $s \in L$, with $|s| \geq p$, s can be divided into 5 pieces $s = uvxyz$ satisfying:

- 1 For each $i \geq 0$, $uv^i \underline{xy}^i z \in L$
- 2 $|vy| > 0$
- 3 $|vxy| \leq p$

- Last week we saw how to use this to prove languages not context-free

The CFG Pumping Lemma

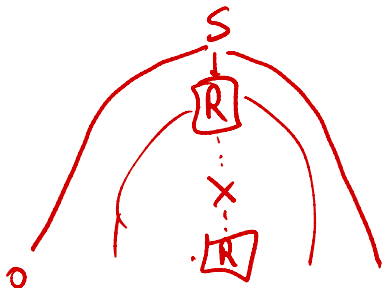
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- But, we did not explain why this lemma is true

Proving the CFG Pumping Lemma (Intuition)



p Rules in our grammar



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 - NFA to DFA using the finger method

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- Understand how to use it.

1. Assume L is reg.

2. Observe that 0^*1^* is reg.



5 Regular Expressions 2. Since reg. lang. are closed under \cap

- Be able to build an RE for a language
- RE to NFA
- NFA to RE

$\Rightarrow L_1$ is regular

6 Regular Language Pumping Lemma

- Remember statement as sequence of quantifiers
- Understand why it is true (state of NFA must repeat)
- Understand how to use it.
- Also know how to prove languages not regular using closure properties

Contradiction!

$$L: \{w \mid \#0's \geq \#1's\} \quad L_1 = \underline{\underline{L \cap 0^*1^*}}$$

$L_1: \{w \mid 0^n 1^n\}$ - not regular

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- 8 Context-free Grammars (CFG)
 - Remember what this means

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- Be able to construct one from language description
- Remember what a derivation is and what a parse tree is
- $\text{PDA} \iff \text{CFG}$ (at a high level)

9 CFL pumping lemma

- There will not be any questions on the CFL pumping lemma on the exam
- But, there will be on the next homework

Exam Format

- 7 questions – most have multiple parts
- Covers most of the material outlined above
- 2 questions requiring proofs, the rest are more constructive
- Some yes/no questions

Don't Forget

- Exam is in class on Thursday 11:10-12:25, don't be late!
- You can bring two 8.5×11 piece of paper

Any Questions?

if L is regular

$\exists p$ s.t. $\forall w \in L$ s.t. $|w| > p$

\exists partition $w \equiv xyz$ s.t.

1) $|y| > 0$

2) $|xy| < p$

3) $\forall i \geq 0 \quad xy^iz \in L$

Any Questions?

Assume L is regular

$\forall p \exists \underline{w} \in L, |w| > p$ s.t.

\forall partitions $w = xyz$

$\exists i$ s.t. $xy^iz \notin L$

Any Questions?

$$L = 0^n 1^{2^n}$$

$$S \rightarrow 0S11 \mid \epsilon$$

Any Questions?

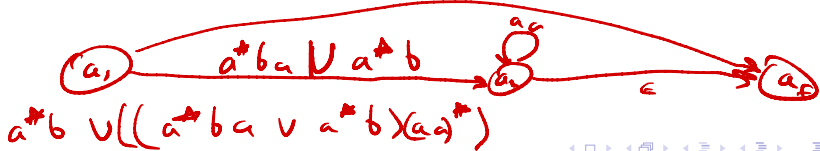
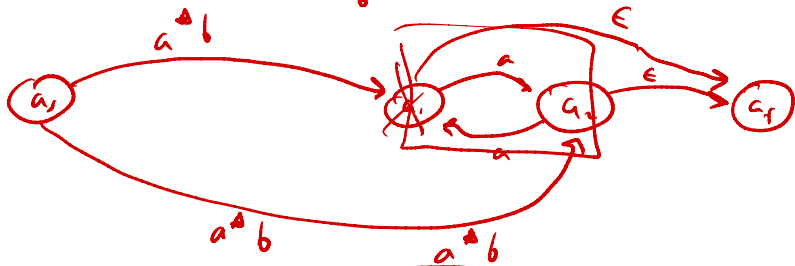
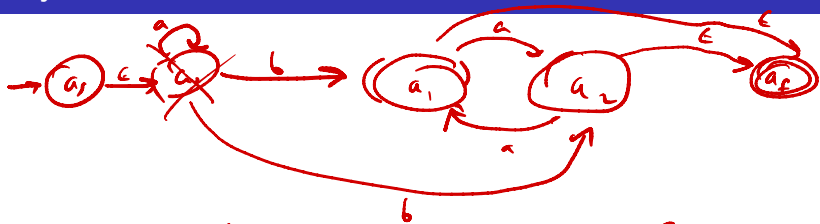
$$L = 0^{n-2} 1^n \quad L_1 = \{000\}$$

$$L_1 = \{w \mid \text{s.t. } w = 0^n 1^n \quad \text{for } n \geq 1\}$$

$$L_1 = 000 \parallel \underline{\underline{0^{n-2} 1^n}}$$

$$L_1 = \underline{\underline{L_2}} \parallel \underline{\underline{L}}$$

Any Questions?



$$a^+b \cup ((a^+b a \cup a^+b)(a)^+)$$